

A Clockless Synchronisation Framework for Cooperating Mobile Robots



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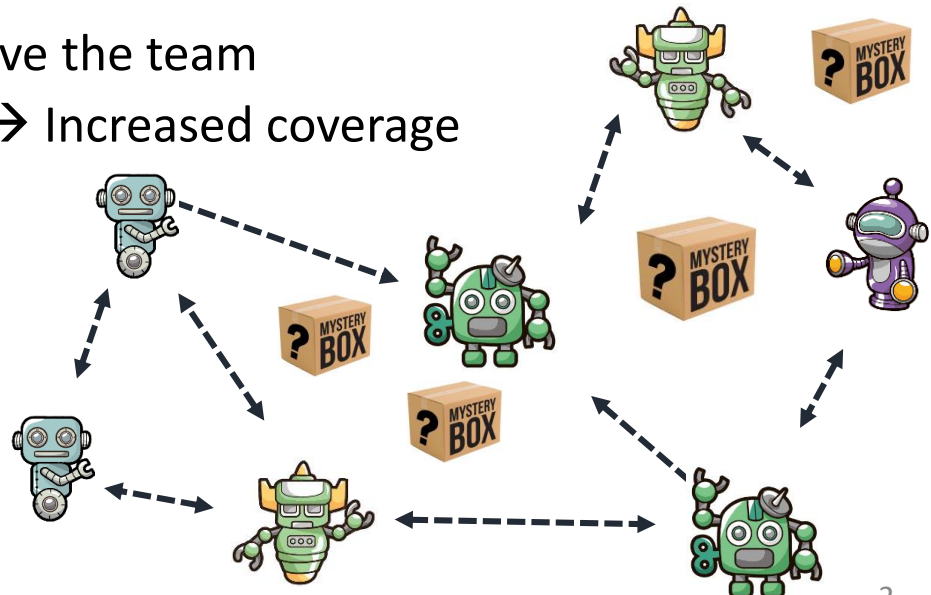
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Motivation

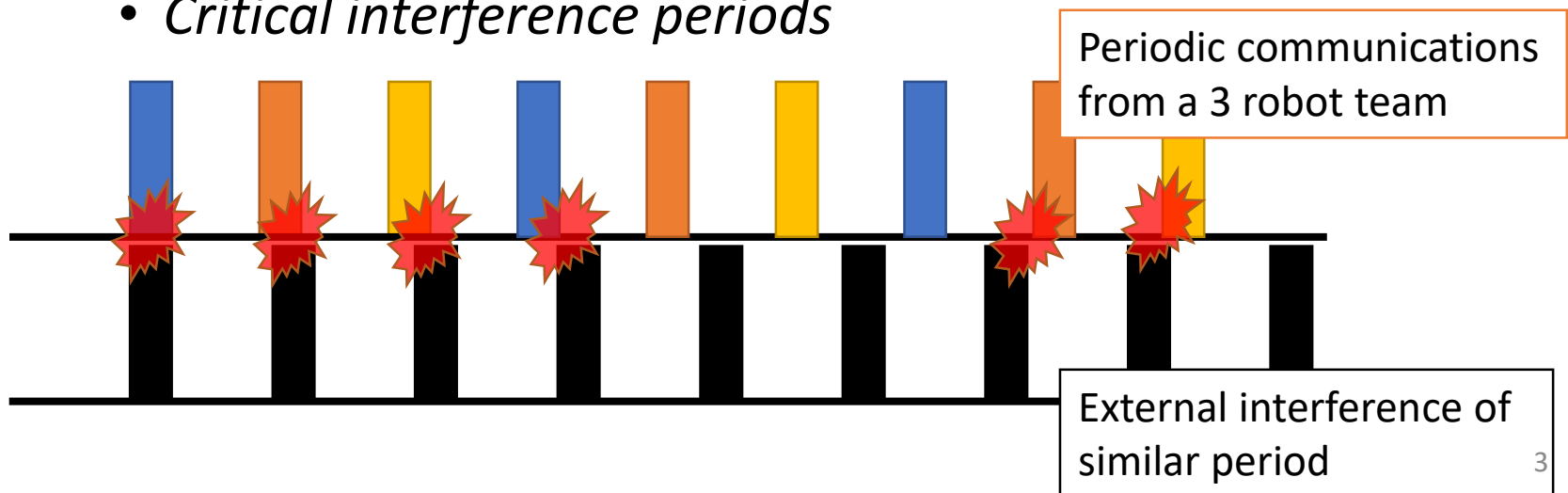
- Widespread use of small teams of cooperative mobile robots
 - Need to communicate using a wireless medium
 - Improved quality of communications → Improved cooperation
 - Are dynamic
 - Robots may join or leave the team
 - Topology can change → Increased coverage



Motivation

In order to improve wireless communications

- Synchronising transmissions is vital
 - Support periodic communications
- Global clock synchronisation
 - Does not cope with periodic external interference
 - *Critical interference periods*



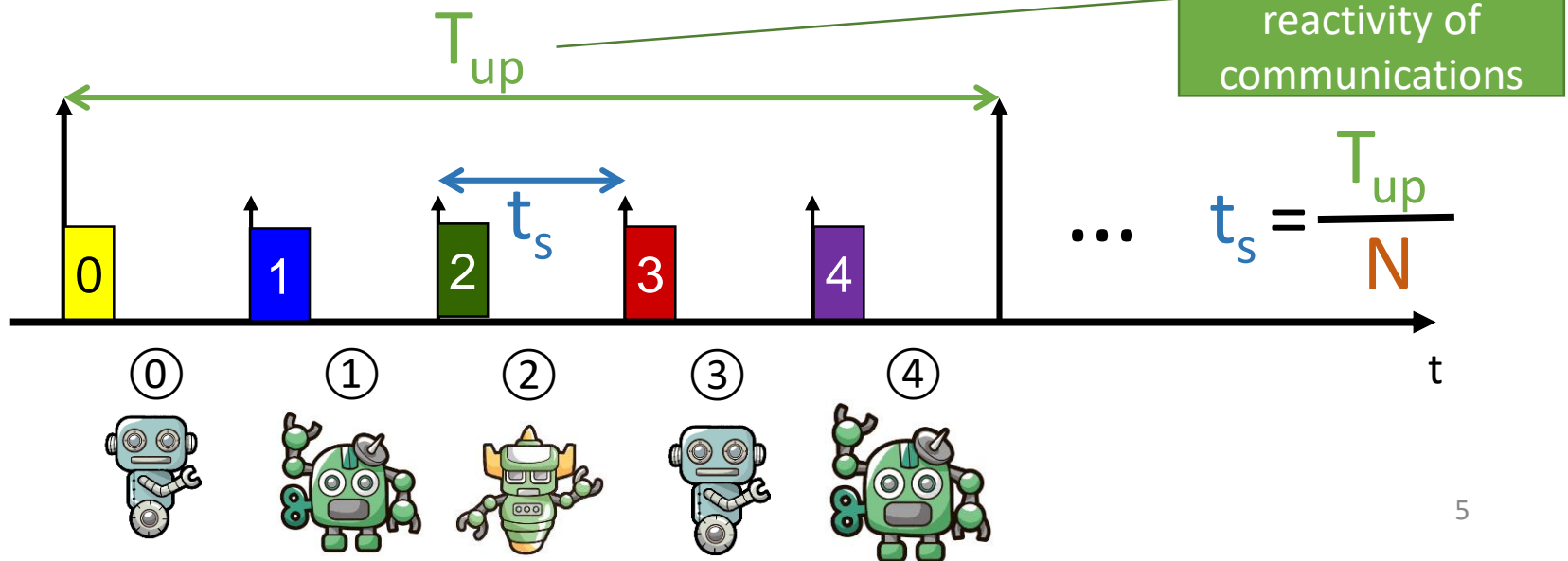
Proposed solution

Use RA-TDMA+ to synchronise transmissions

What is RA-TDMA+ ?

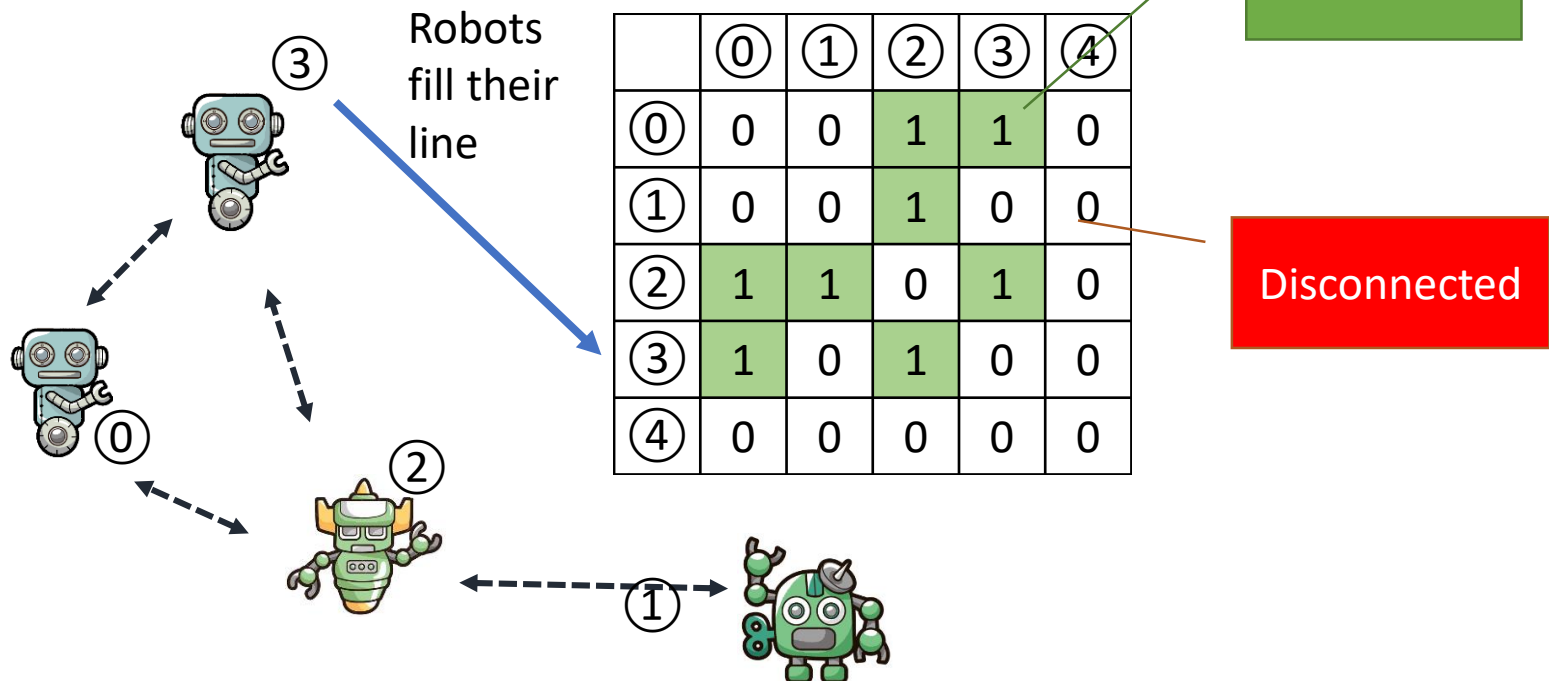
Overlay communication protocol (uses CSMA/CS)

- Creates a TDMA communications round
 - Fixed round duration \rightarrow UPdate period (T_{up})
 - Divides in N slots \rightarrow One slot dedicated to each robot
 - No spatial reutilization (Robots are mobile!)



Dynamic team topology

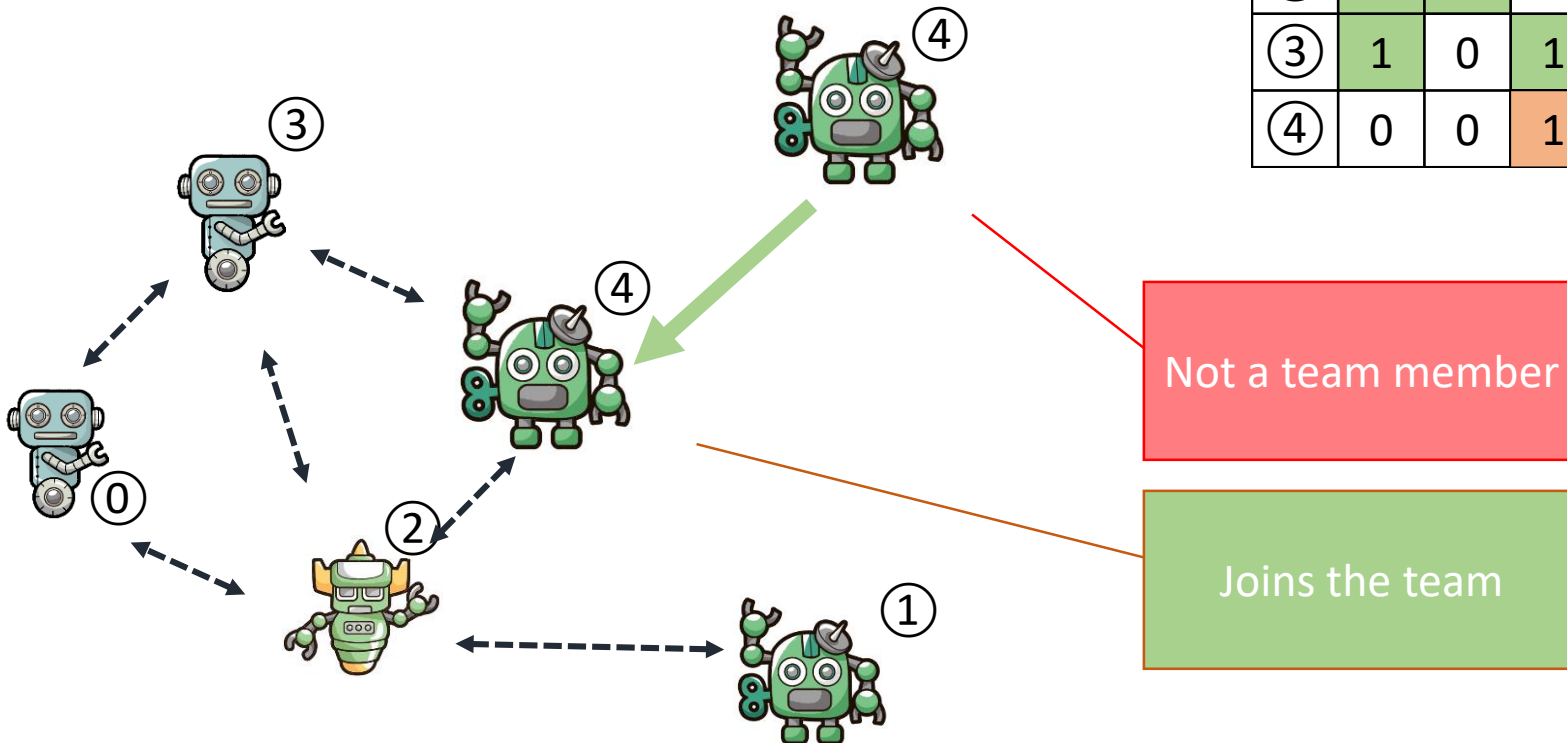
- Supports dynamic topologies (mesh networks)
 - Team topology may change on-line
 - Connectivity matrix is disseminated
 - Filled cooperatively



Dynamic team membership

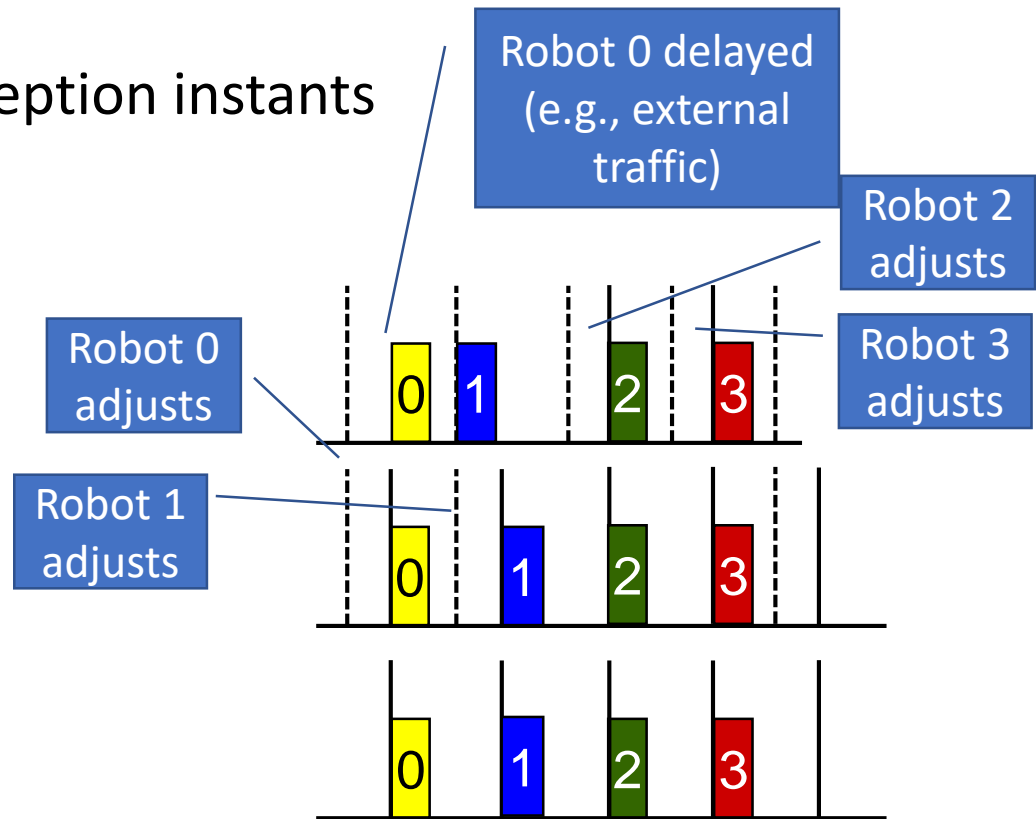
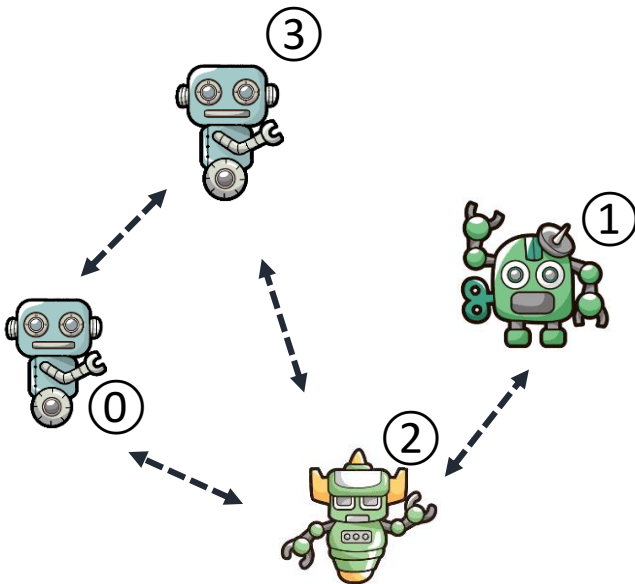
- Dynamic teams
 - Team composition may change on-line
 - Slots are created/destroyed dynamically

	①	②	③	④	
①	0	0	1	1	0
②	0	0	1	0	0
③	1	1	0	1	1
④	1	0	1	0	1
⑤	0	0	1	1	0



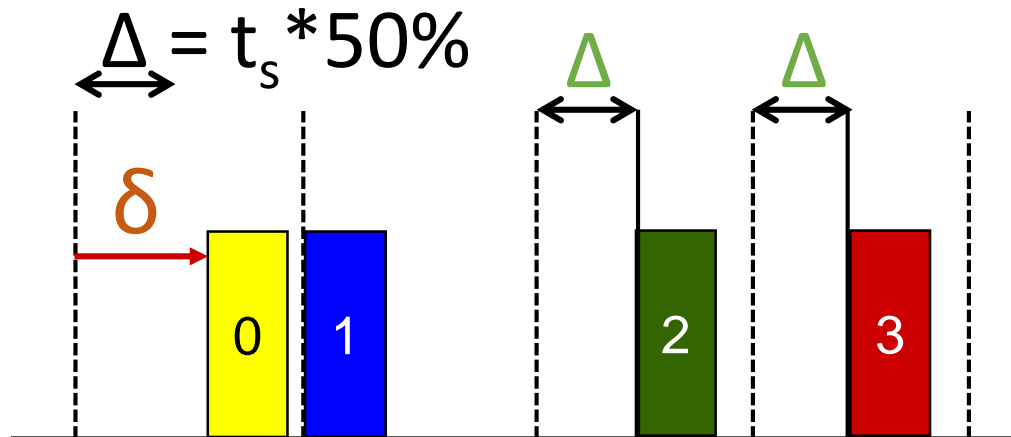
Creating and maintaining the TDMA

- RA-TDMA+ is a loose synchronisation protocol
 - No global clock
 - Uses message reception instants



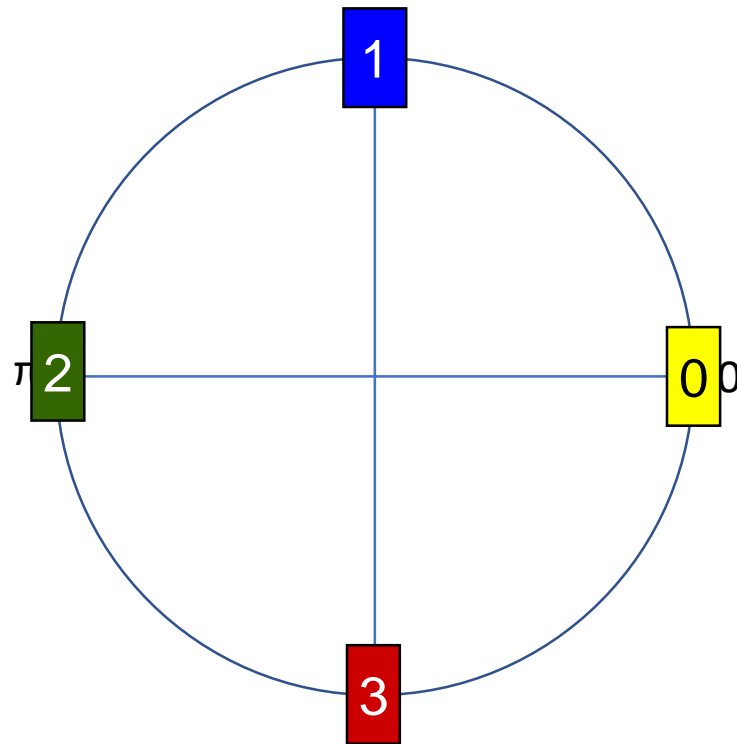
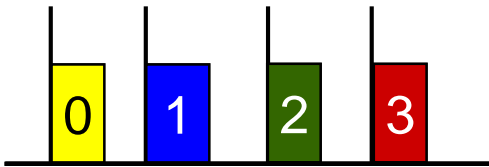
The synchronisation protocol

- Robots **ONLY** delay their communications
 - Anticipating leads to increased BW
 - Parameter Δ limits the maximum delay per round



Reframing the problem

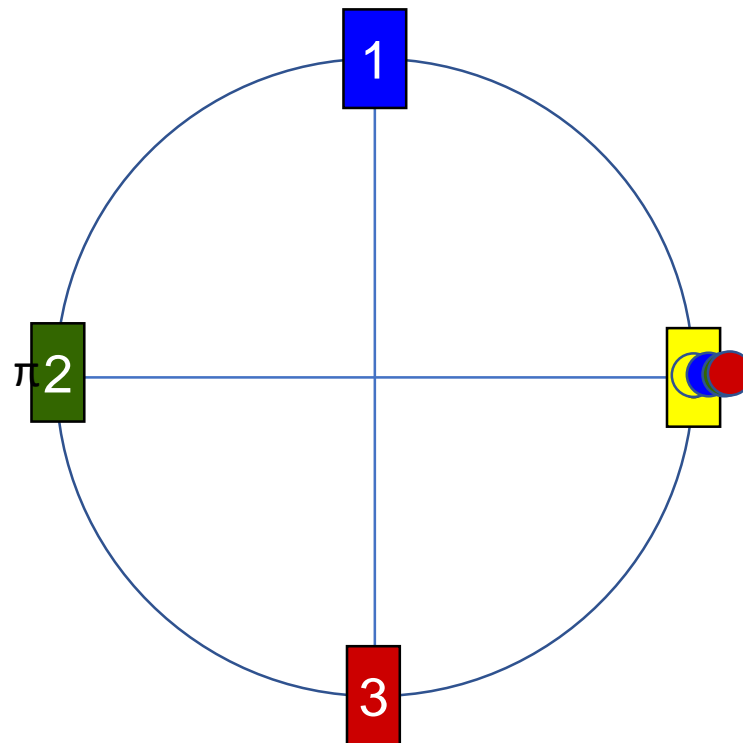
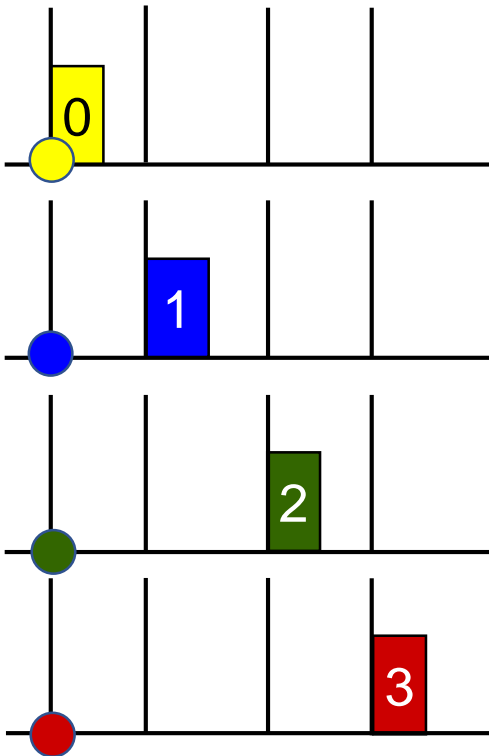
- TDMA round is cyclic
 - We can represent round and slot time as a phase



Reframing the problem

Using simple maths

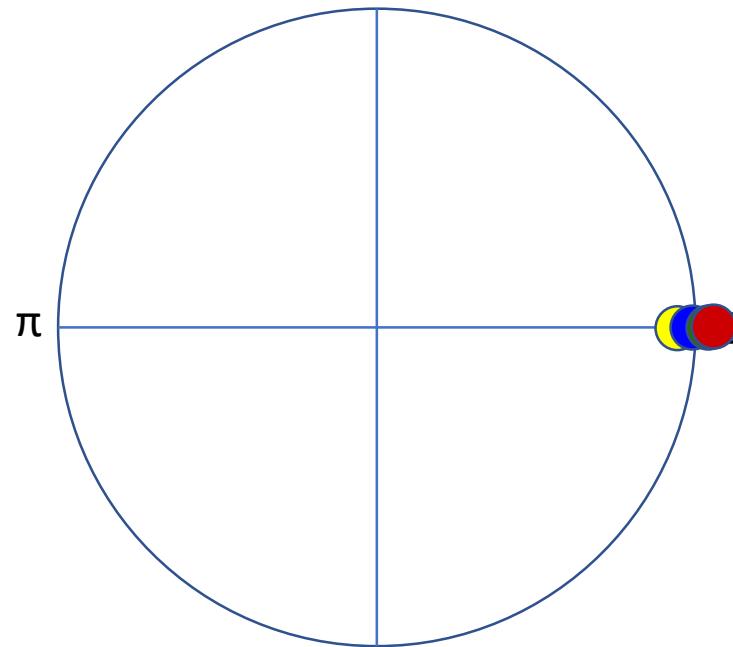
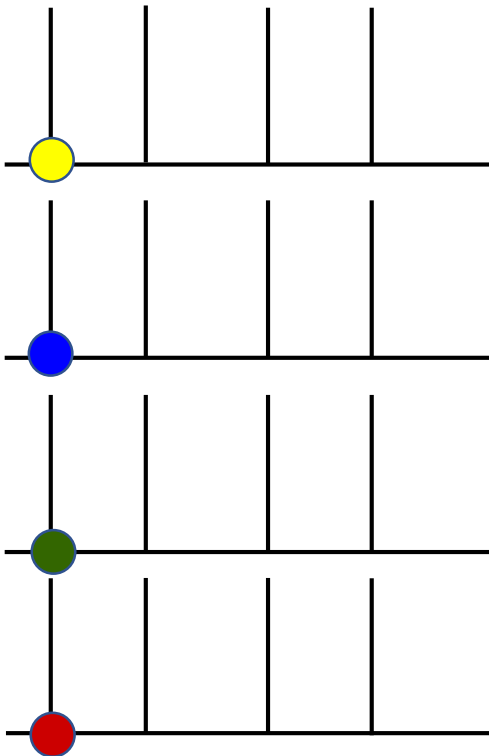
- Message arrival times \rightarrow round start phases (Φ)



Reframing the problem

Making adjustments to the round (delaying comms.)

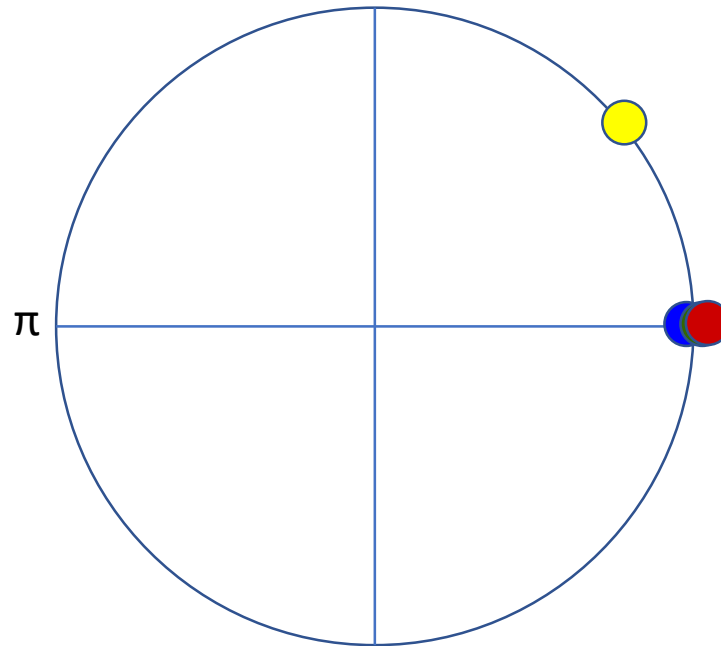
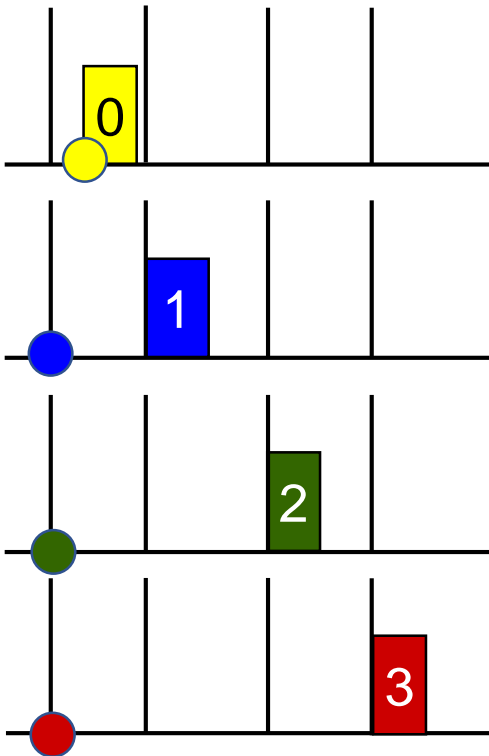
- Makes initial round phases increase



Reframing the problem

Problem can now be stated as:

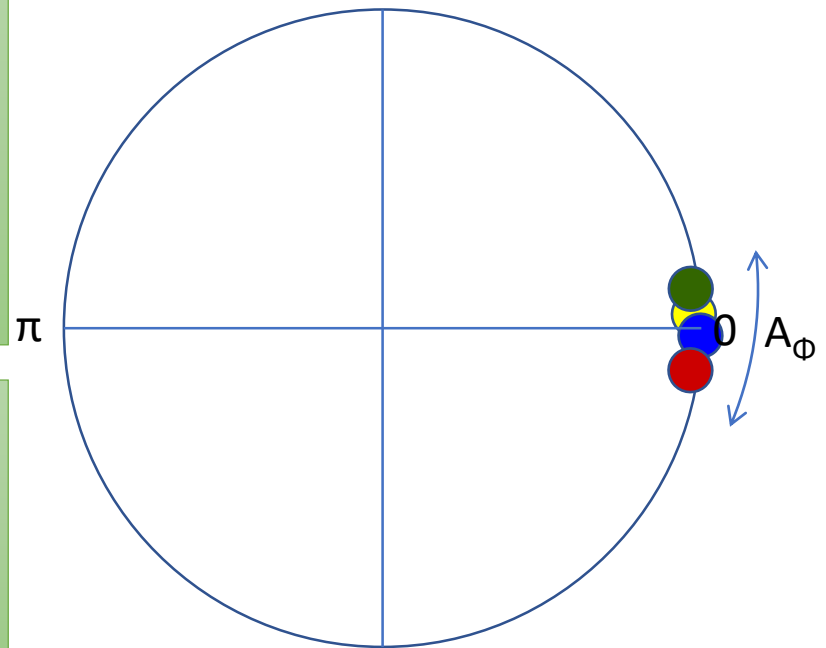
- Advance the round start phase until it matches the most advanced



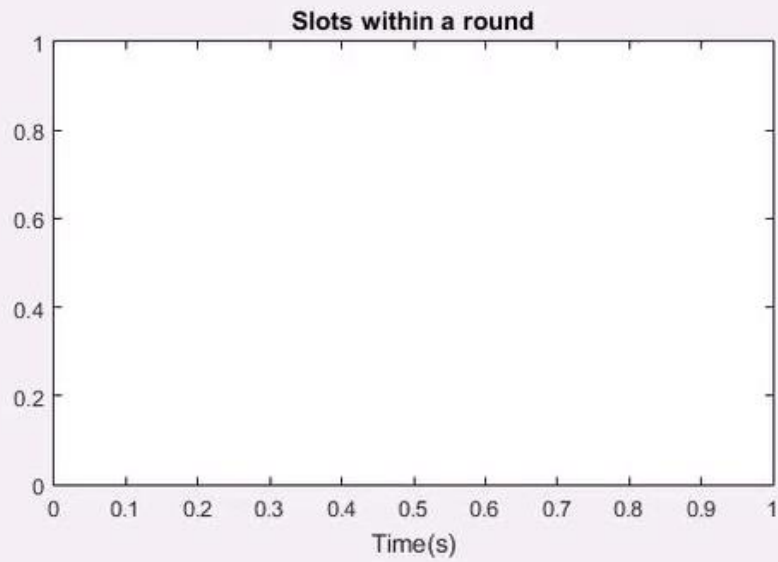
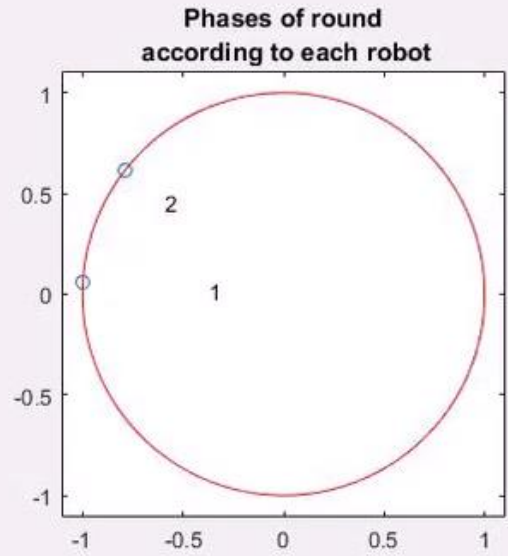
Reframing the problem

Robots are synchronised when there is an agreement on the round start phase

Or equivalently
When the largest difference between phases (A_Φ) is zero



Video



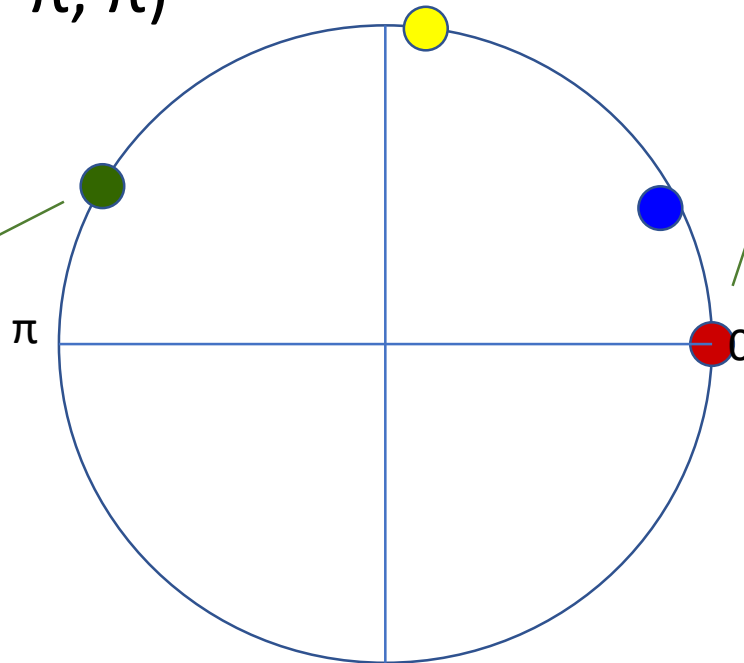
When does it work?

If $A_\Phi < \pi$

And we consider Φ in $[-\pi, \pi)$

Independently from the point of observation

Then the green Φ is the most advanced



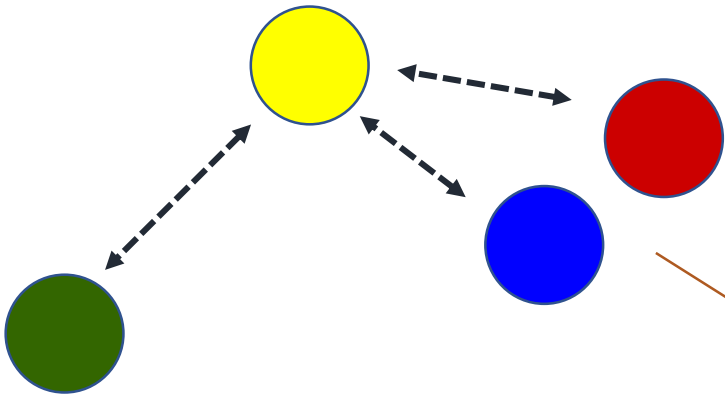
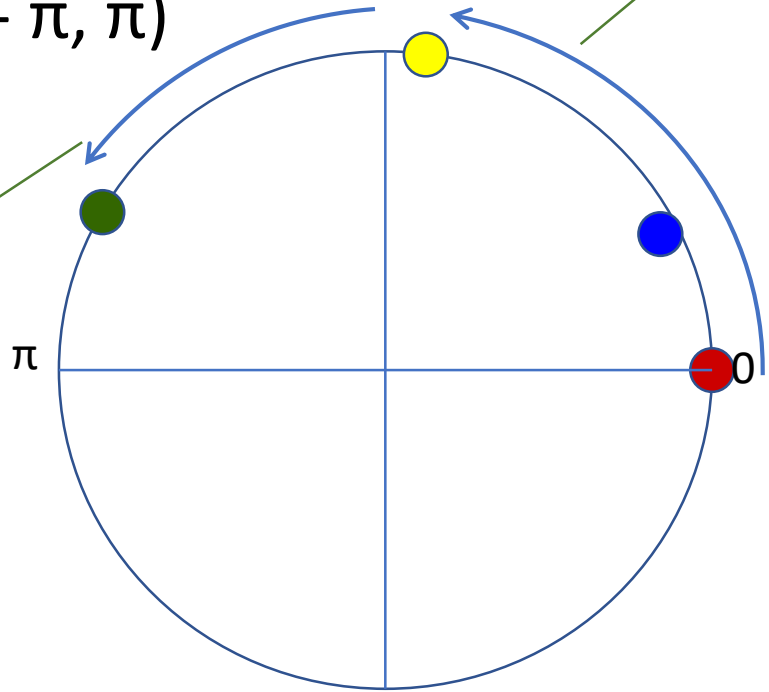
When does it work?

If $A_\Phi < \pi$

And we consider Φ in $[-\pi, \pi)$

Or towards a phase that is moving towards the most advanced

Phases will either advance towards the most advanced

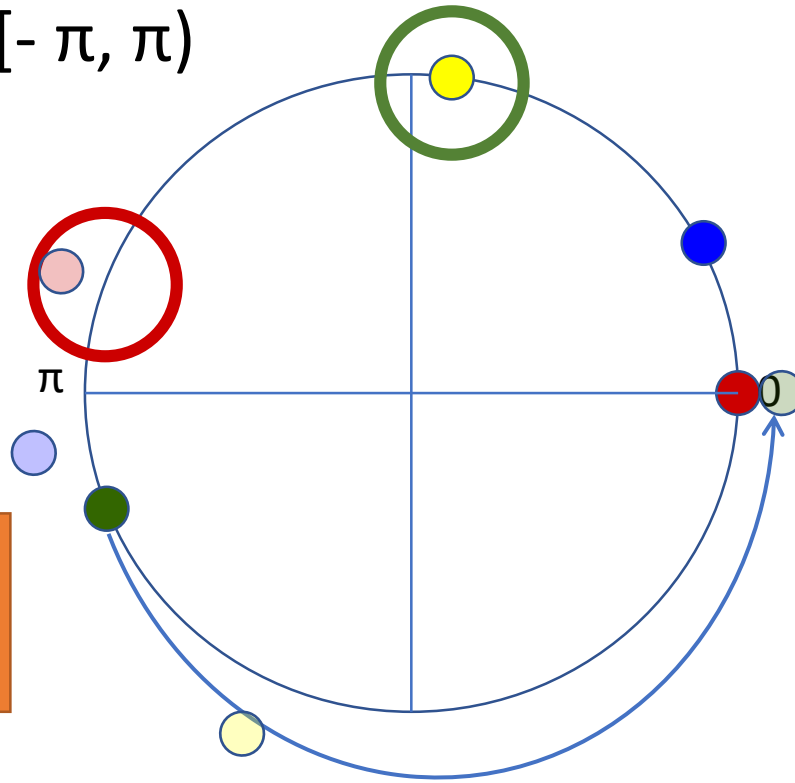


The topology is not (necessarily) fully connected

When it may fail?

If $A_\Phi \geq \pi$

And we consider Φ in $[-\pi, \pi)$



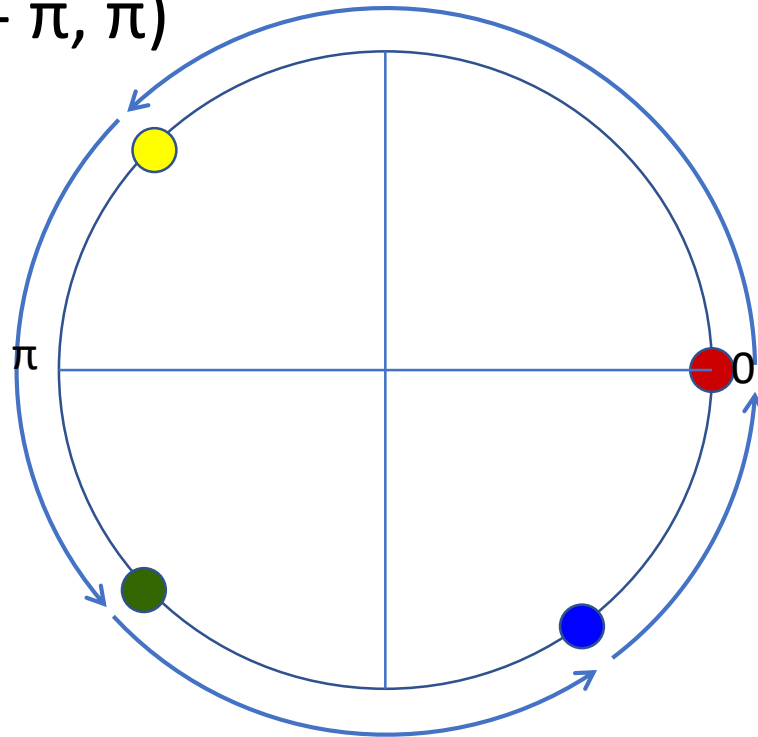
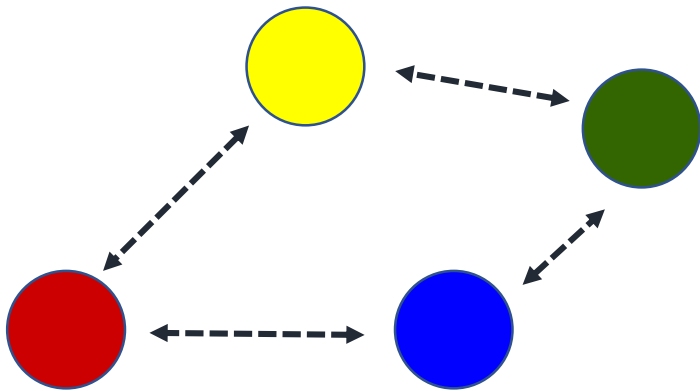
The most advanced Φ depends on the point of observation

When it may fail?

If $A_\Phi \geq \pi$

And we consider Φ in $[-\pi, \pi)$

Phases may chase each other continually for large Δ

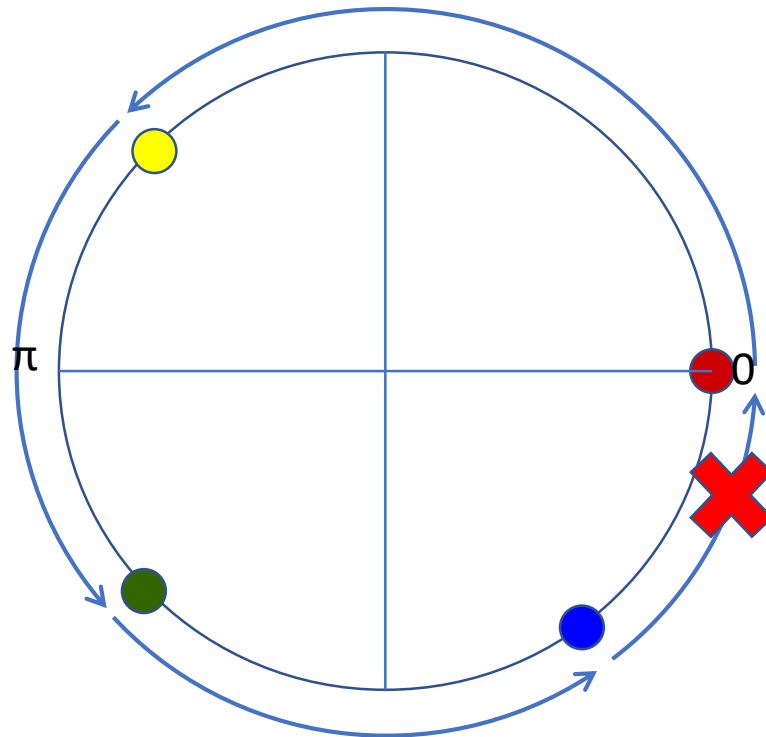
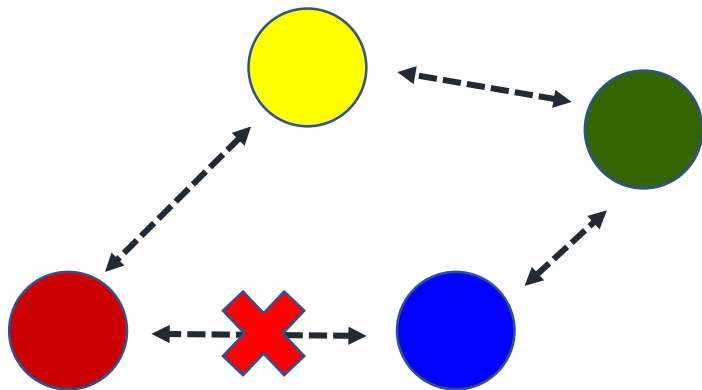


Heuristic to address $A_{\phi} \geq \pi$

Solution:
Break the cycle

Using the topology information
on the connectivity matrix

Create a spanning tree
(for synchronisation ONLY)



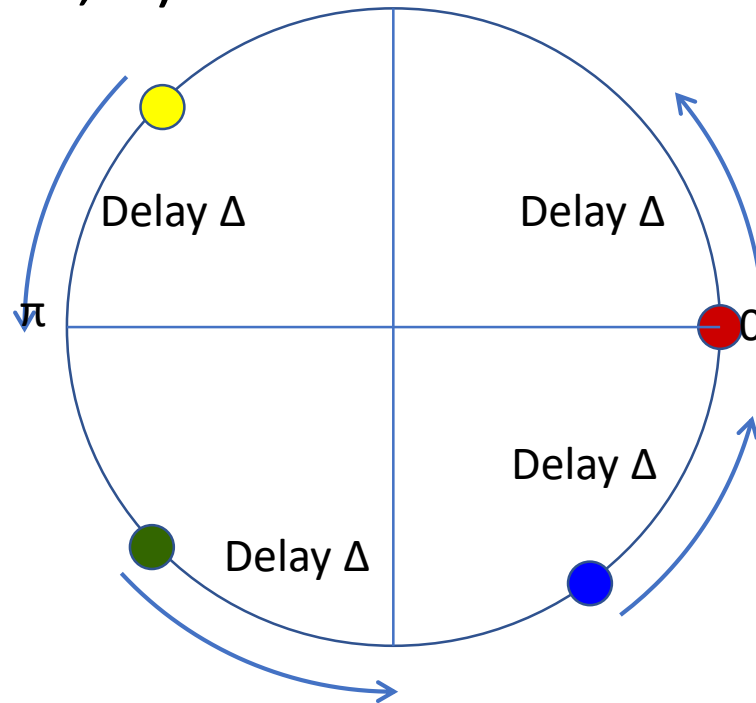
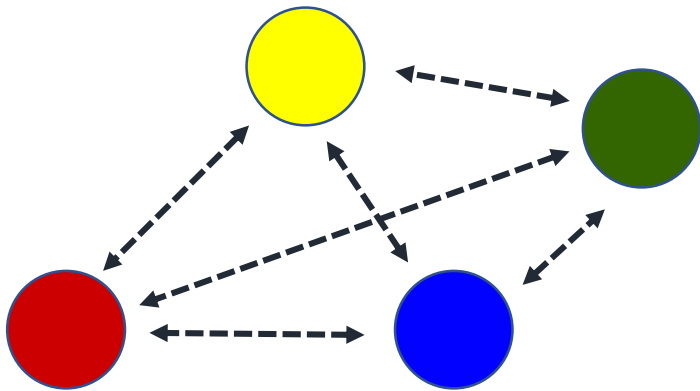
Drawbacks!!!
Longer time to propagate

When it may fail?

If $A_\Phi \geq \pi$

And we consider Φ in $[-\pi, \pi)$

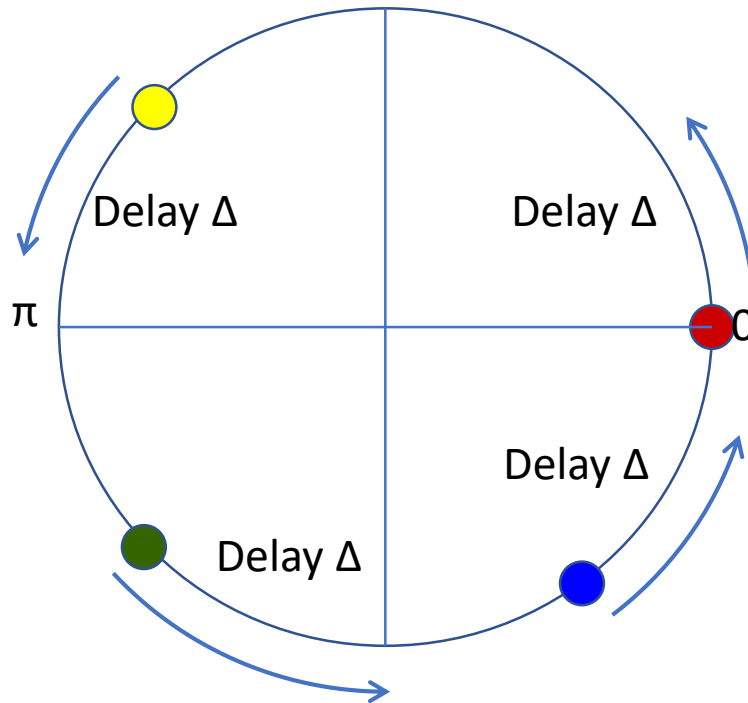
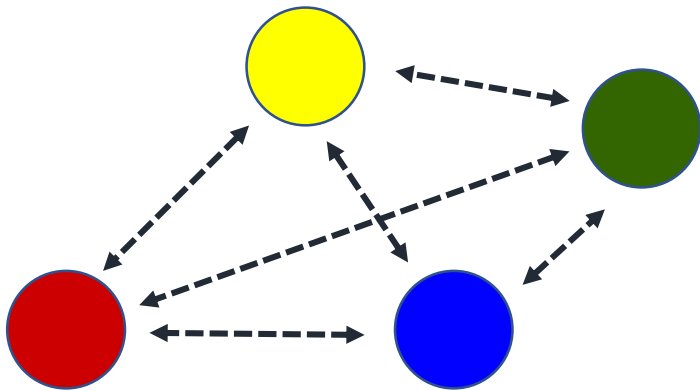
Phases may all delay the same amount every round
(Δ – The maximum delay)



Heuristic to address $A_{\phi} \geq \pi$

Solution:
Randomise Δ

E.g., $\Delta' = \Delta(0.8 + 0.2 * \text{rnd})$



Experimental results

The following experimental results measure

- Impact of different Δ values
- Ability to maintain synchronisation
 - under different load conditions
- Ability to maintain synchronisation
 - Varying team composition and topology

Matlab simulations

Experimental results – Impact of Δ

$$A_{\Phi} < \pi$$

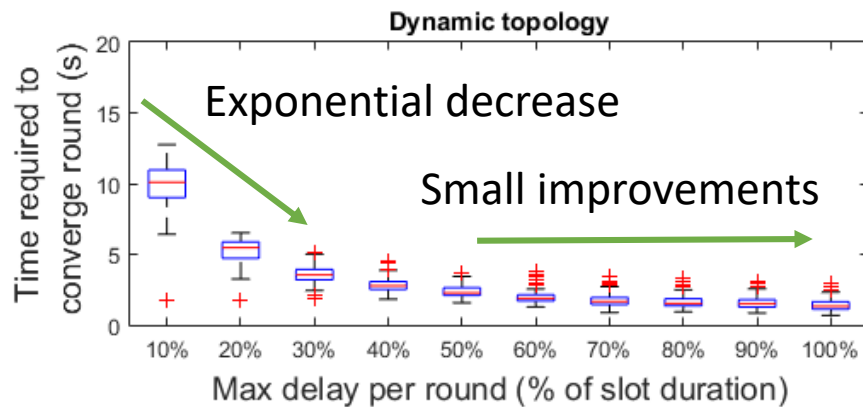
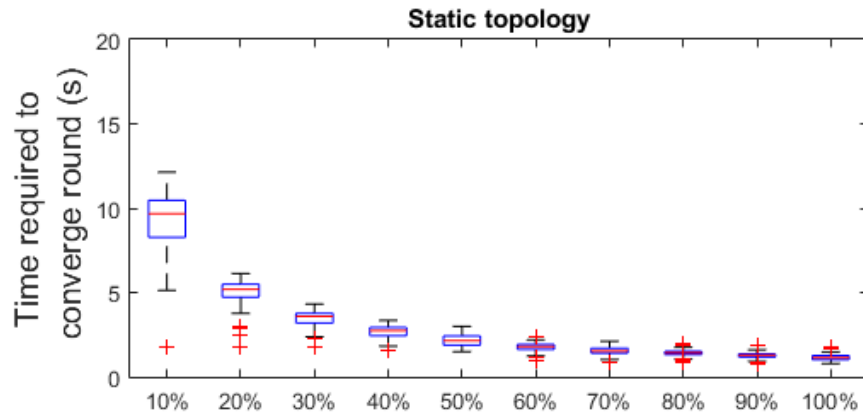
Normal operation
Small delays introduced by external traffic, or OS induced

10 robots
 $T_{up} = 200\text{ms}$

Random starting phases Φ

Negligible impact of movement

All robots synchronised in every case



Changing positions every 10s

Experimental results – Impact of Δ

Exceptional conditions
Team membership changes

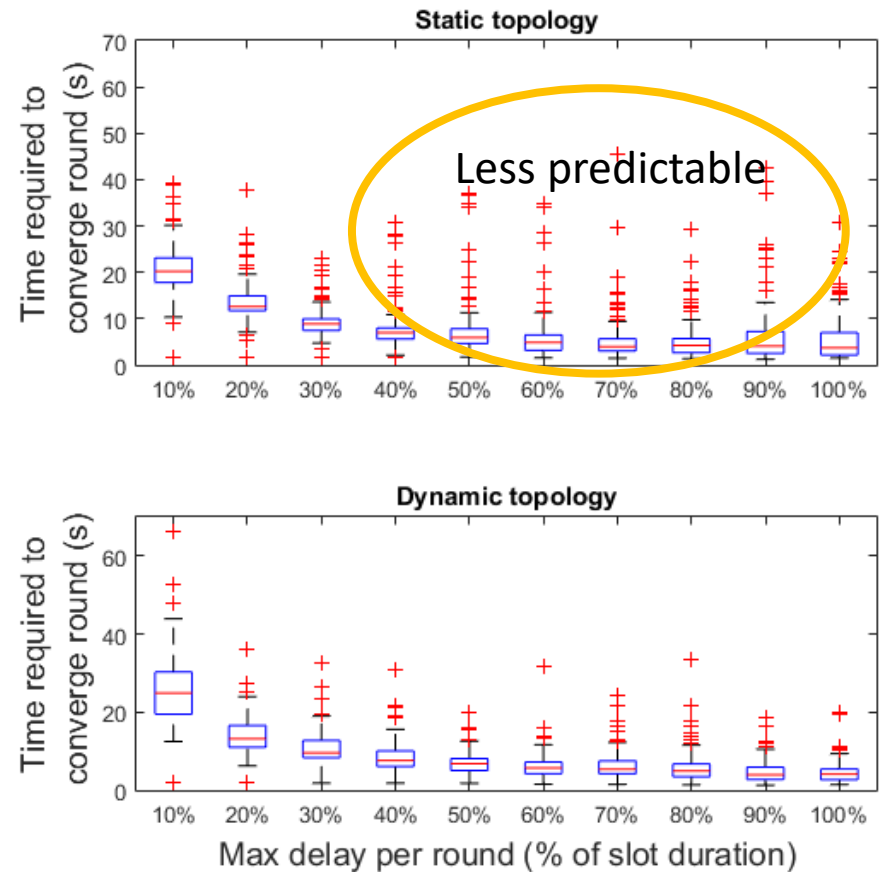
10 robots
 $T_{up} = 200\text{ms}$

Random starting phases Φ

Negligible impact of movement

All robots synchronised in
every case

$$A_{\Phi} \geq \pi$$



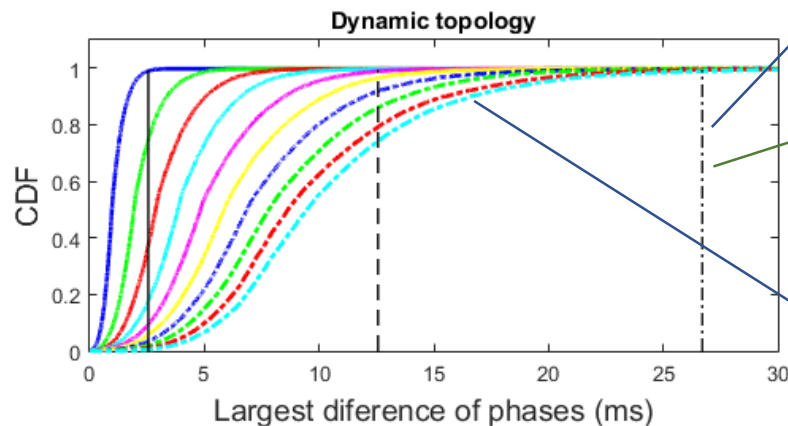
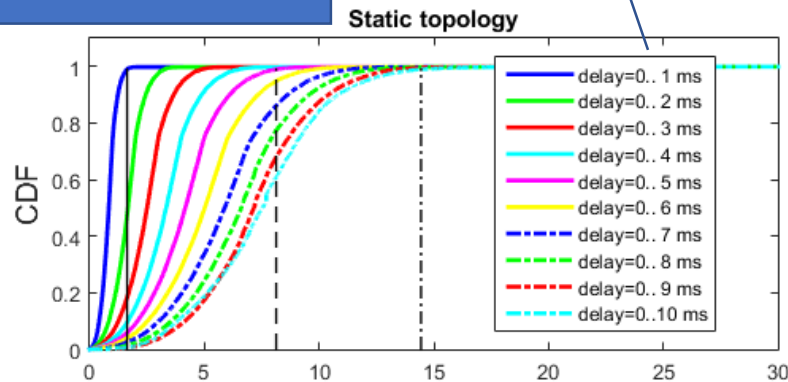
Changing positions every 10s

Experimental results

Impact of delays

$\Delta=40\%$ of slot duration

Injected uniform random delays



$A_{\Phi} < \pi$ (<100ms)
Due to no membership changes

Graceful degradation!

99 percentile for 1,5, and 10ms

Shifted to the right with dynamic topology (<30ms)

Measured the maximum difference of phases for each maximum delay

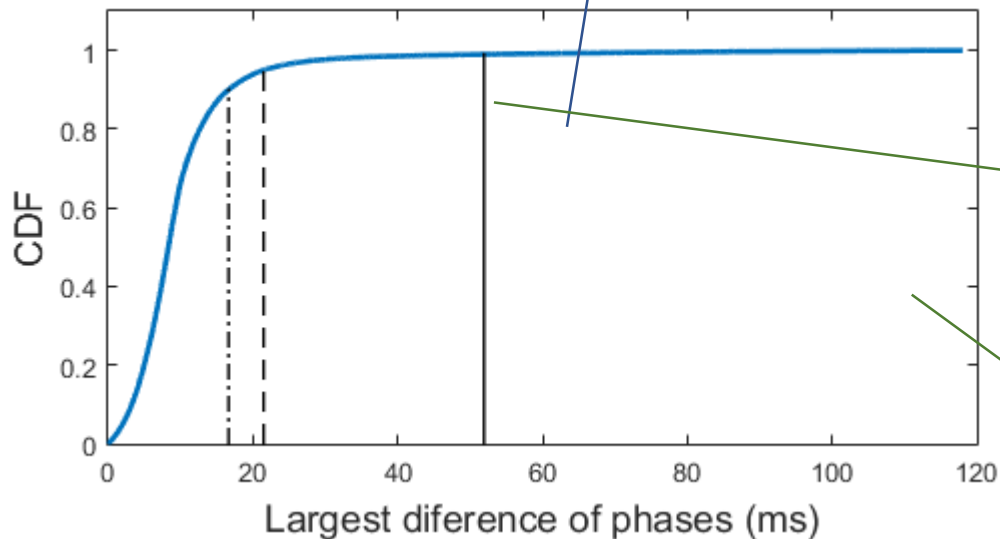
Experimental results

Changing team membership

Number of robots vary from 3 \rightarrow 10 \rightarrow 3 ...
Add/remove robot every 50s for 10000s
(200 team membership changes)

$\Delta=40\%$ of slot
duration

Random delays
Uniform 0..10ms



99% $\ll \pi$
(100ms)

$> \pi$
When team changes

Conclusions

Proposed a mechanism to create a global TDMA

- Fully distributed
- Dynamic robot teams
- No global clock
- In normal operation → convergence guaranteed
- Exceptional conditions → heuristic was proposed

Questions?